

# Linear Parallel Interference Cancellation Using Fixed Weighting Factors for Long-Code CDMA

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**Abstract** — Linear weighted multistage parallel interference cancellation (PIC) implements exactly the family of polynomial expansion detectors. For long-code CDMA, a set of optimal weights is found which minimizes the ensemble averaged mean squared error (MSE) over random codes. The weights are dependent on moments of the eigenvalues of the correlation matrix, where exact expressions are derived. The loss incurred by averaging rather than using the optimal, time-varying weights is practically negligible.

## I. INTRODUCTION

Consider a  $K$ -user symbol-synchronous CDMA system with processing gain  $N$ . The received signal vector is  $\underline{r} = \underline{A}\underline{d} + \underline{n}$ , where  $\underline{A} = (\underline{a}_1, \underline{a}_2, \dots, \underline{a}_K)$  is the matrix containing all users' spreading codes,  $\underline{d} = (d_1, d_2, \dots, d_K)^T$  the data vector and  $\underline{n}$  the AWGN with variance  $\sigma^2$ .

The detailed structure of the  $i$ th PIC stage with weight  $\mu_i$  is depicted here where MF denotes matched filtering. A multistage PIC is a simple cascade of  $m$  PIC stages, whose output is

$$\underline{y}_m = \left[ \underline{I} - \prod_{i=1}^m (\underline{I} - \mu_i (\underline{R} + \sigma^2 \underline{I})) \right] (\underline{R} + \sigma^2 \underline{I})^{-1} \underline{A}^H \underline{r} \quad (1)$$

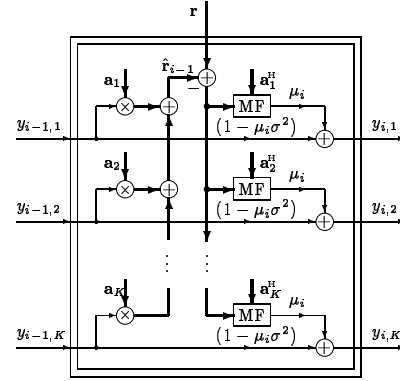
where  $\underline{R} = \underline{A}^H \underline{A}$  is the correlation matrix. By choosing an appropriate set of weights, the PIC can implement exactly any detector of the form of a polynomial in  $\underline{R}$  applied to the code matched-filtered output  $\underline{A}^H \underline{r}$ .

It has previously been shown that the PIC is a realization of the steepest descent algorithm used to minimize the MSE. Following this interpretation, a unique set of weights, dependent on the eigenvalues of  $\underline{R}$ , was found to lead to the minimum achievable MSE for a given number of stages in a short-code system [1]. This approach is too complex for long-code systems. Instead, we consider using a set of code-invariant weights designed to give the minimum ensemble averaged MSE over random codes.

The ensemble average of the excess MSE, as compared to the MMSE, is expressed as a function of  $\underline{u} = (\mu_1, \mu_2, \dots, \mu_m)$

$$J^{(m)}(\underline{u}) = E \left\{ \sum_{k=1}^K \frac{\lambda_k}{\lambda_k + \sigma^2} \prod_{i=1}^m |1 - \mu_i (\lambda_k + \sigma^2)|^2 \right\} \quad (2)$$

where the  $\lambda_k$ 's are the eigenvalues of  $\underline{R}$  and the expectation is taken over random codes. By an elementary symmetric polynomial transform in  $\underline{u}$  we can rewrite (2) as a quadratic function in a vector  $\underline{z}$ , which is a function of  $\underline{u}$ . A unique



minimum is then obtained where the corresponding  $\mu_k$ 's are found as the inverse polynomial transform.

For an  $m$ -stage PIC, the weights depend on the first  $2m$  moments of the eigenvalues, defined as  $M_r = E\{\lambda^r\}$ ,  $r = 1, 2, \dots, 2m$  where  $\lambda$  is an arbitrary eigenvalue of  $\underline{R}$ . Moreover,

$$\begin{aligned} M_r &= \frac{1}{K} E \{ \text{tr} \{ \underline{R}^r \} \} \\ &= \frac{1}{K} \sum_{k_1=1}^K \sum_{k_2=1}^K \cdots \sum_{k_r=1}^K E \{ R_{k_1 k_2} R_{k_2 k_3} \cdots R_{k_{r-1} k_r} R_{k_r k_1} \} \\ &= \frac{1}{K} \sum_{k_1=1}^K \sum_{k_2=1}^K \cdots \sum_{k_r=1}^K \sum_{n_1=1}^N \sum_{n_2=1}^N \cdots \sum_{n_r=1}^N \\ &\quad E \{ A_{n_1 k_1}^* A_{n_1 k_2} A_{n_2 k_2}^* A_{n_2 k_3} A_{n_3 k_3}^* \cdots A_{n_r k_r}^* A_{n_r k_1} \}. \quad (3) \end{aligned}$$

Since  $A_{nk}$  are all independent random variables, only terms containing all complex conjugate pairs are relevant.  $M_r$  is then obtained through evaluation of the summation over all combinations of indices. As the expectation is taken over all code-sets,  $M_r$  only depends on  $N$  and  $K$ , but not on specific codes. In fact it is a polynomial in  $N$  and  $K$  [2].

With the exact expressions derived, the moments can be evaluated easily and the optimal weights computed. The computational complexity is minor for a moderate number of stages and hence can be implemented on-line. Simulation results show that the penalty of averaging rather than using the optimal weights dependent on the instantaneous spreading codes is negligible in most cases of interest.

## REFERENCES

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