

# **Network Game with a Probabilistic Description of User Types**

**Hongxia Shen and Tamer Başar**

**Coordinated Science Laboratory**

**University of Illinois at Urbana-Champaign**

**`hshen1, tbasar@control.csl.uiuc.edu`**

**CDC 2004, Atlantis, Bahamas**

**December 17, 2004**

# Outline

- Previous Work
- Problem Formulation
  - Complete Information
  - Partially Incomplete Information
  - Totally Incomplete Information
- Single User
- Two Users
- Multiple Users
- Conclusions and Extensions

## Previous Work

➤ Başar and Srikant (2002a, 2002b)

Hierarchical Stackelberg Network Game Model

Solution under Uniform Pricing

Asymptotic Behavior Analysis

➤ Shen and Başar (2004)

Solution under Differentiated Pricing

Comparison with Uniform Pricing and Asymptotic Behavior Analysis

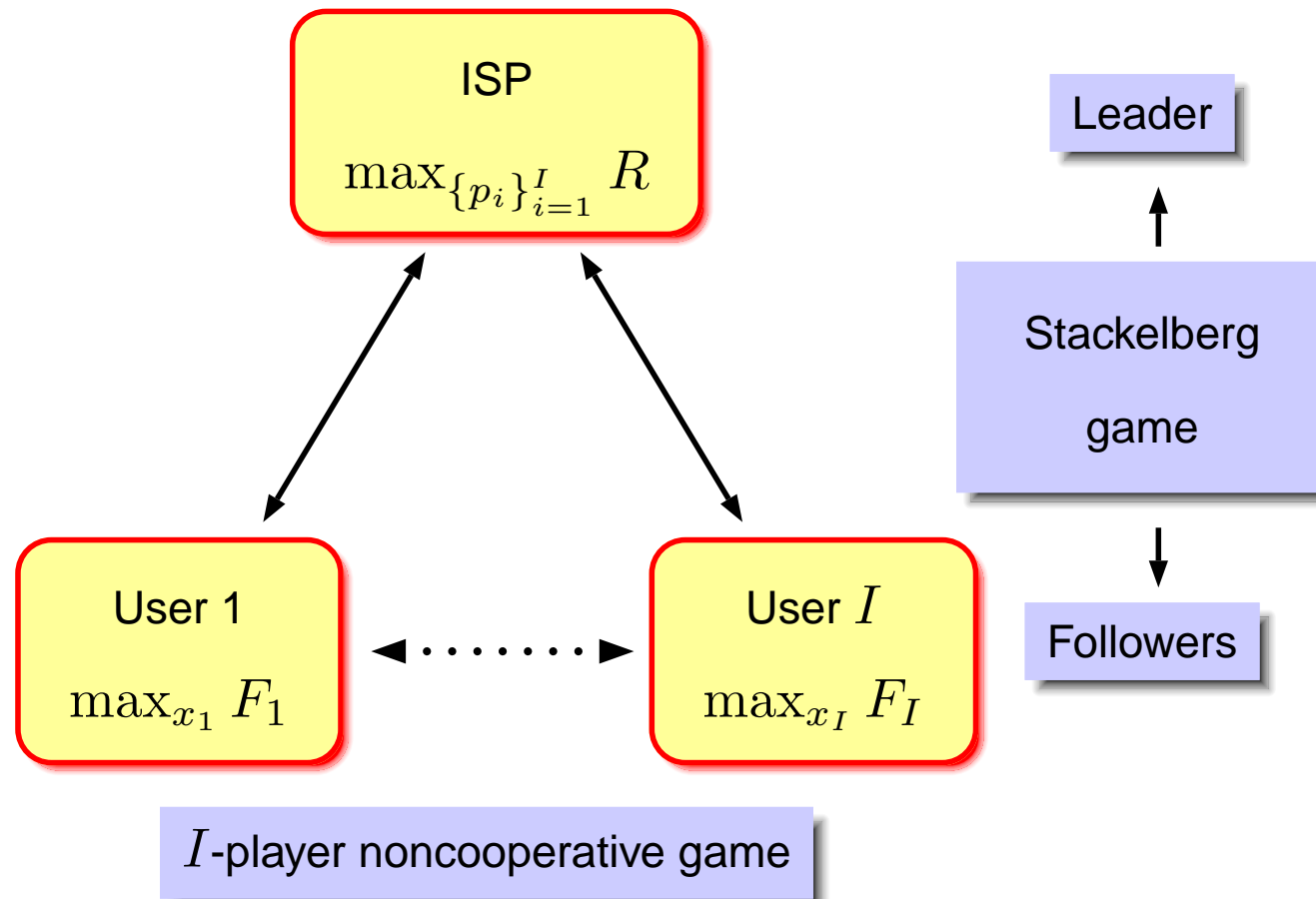
---

T. Başar and R. Srikant (2002a), “Revenue-maximizing pricing and capacity expansion in a many-users regime,” *Proc. IEEE INFOCOM 2002*, pp. 1556–1563.

T. Başar and R. Srikant (2002b), “A Stackelberg network game with a large number of followers,” *J. Optimization Theory and Applications*, 115(3): 479–490.

H.-X. Shen and T. Başar (2004), “Differentiated Internet pricing using a hierarchical network game model,” *Proc. IEEE ACC 2004*, pp. 2322–2327.

## Two-Level Hierarchical Network Game



## Problem Formulation

- Single Internet Service Provider (ISP)
- Set of users:  $\mathcal{I} = \{i : 1 \leq i \leq I\}$
- Flow vector:  $\vec{x} = (x_1, \dots, x_I)^T \in \Omega$
- Price vector:  $\vec{p} = (p_1, \dots, p_I)^T$
- User  $i$ 's net utility:  $f_i(x_i) + g(\vec{x}) - p_i x_i$
- ISP's profit:  $\vec{p}^T \vec{x}$
- Classification of games based on the information structure:
  - ① Complete information
  - ② Partially incomplete information
  - ③ Totally incomplete information

## Complete Information

- The utility function of each user is common knowledge for all the users and the ISP.
- Given  $\vec{p}$ , a Nash equilibrium is  $\vec{x}(\vec{p}) = (x_1(\vec{p}), \dots, x_I(\vec{p}))^T$ , such that for User  $i$ ,

$$x_i(\vec{p}) = \arg \max_{x_i} f_i(x_i) + g(x_i, x_{-i}(\vec{p})) - p_i x_i,$$

where  $x_{-i} = \{x_j\}_{j \neq i}$ .

- There exists a unique Nash equilibrium. See Başar and Srikant (2002a).
- The optimal price vector for the ISP is:

$$\vec{p}^* = \arg \max_{\vec{p}: \vec{p} \geq \theta} \sum_{i \in \mathcal{I}} p_i x_i(\vec{p}).$$

---

T. Başar and R. Srikant (2002a), “Revenue-maximizing pricing and capacity expansion in a many-users regime,” *Proc. IEEE INFOCOM 2002*, pp. 1556–1563.

## Partially Incomplete Information

- The utility function of each user is common knowledge for all the users, but not for the ISP.
- The distribution of the users' types,  $\vec{w} = (w_1, \dots, w_I)^T$ , that determine  $f_i$ 's ( $g$  is deterministic), is known to the ISP.
- Given  $\vec{p}$ , for each fixed  $\vec{w}$ , the unique Nash equilibrium is  $\vec{x}^{\vec{w}}(\vec{p}) = (x_1^{\vec{w}}(\vec{p}), \dots, x_I^{\vec{w}}(\vec{p}))^T$ , such that for User  $i$ ,

$$x_i^{\vec{w}}(\vec{p}) = \arg \max_{x_i} f_i^{w_i}(x_i) + g(x_i, x_{-i}^{\vec{w}}(\vec{p})) - p_i x_i.$$

- The optimal price vector for the ISP is:

$$\vec{p}^* = \arg \max_{\vec{p}: \vec{p} \geq \theta} \sum_{i \in \mathcal{I}} p_i E_{\vec{w}}[x_i^{\vec{w}}(\vec{p})].$$

## Totally Incomplete Information

- User  $i$  has private information  $w_i$  and knows the distribution of  $w_{-i} = \{w_j\}_{j \neq i}$ .
- The distribution of  $\vec{w} = (w_1, \dots, w_I)^T$  is known to the ISP.
- Given  $\vec{p}$ , the pure strategy Bayesian equilibrium, if it exists, is  $x_i^{w_i}(\vec{p})$  for User  $i$ , such that

$$x_i^{w_i}(\vec{p}) = \arg \max_{x_i} f_i^{w_i}(x_i) + E_{w_{-i}}[g(x_i, x_{-i}^{w_{-i}}(\vec{p}))] - p_i x_i.$$

- The optimal price vector for the ISP is:

$$\vec{p}^* = \arg \max_{\vec{p}: \vec{p} \geq \theta} \sum_{i \in \mathcal{I}} p_i E_{\vec{w}}[x_i^{w_i}(\vec{p})].$$

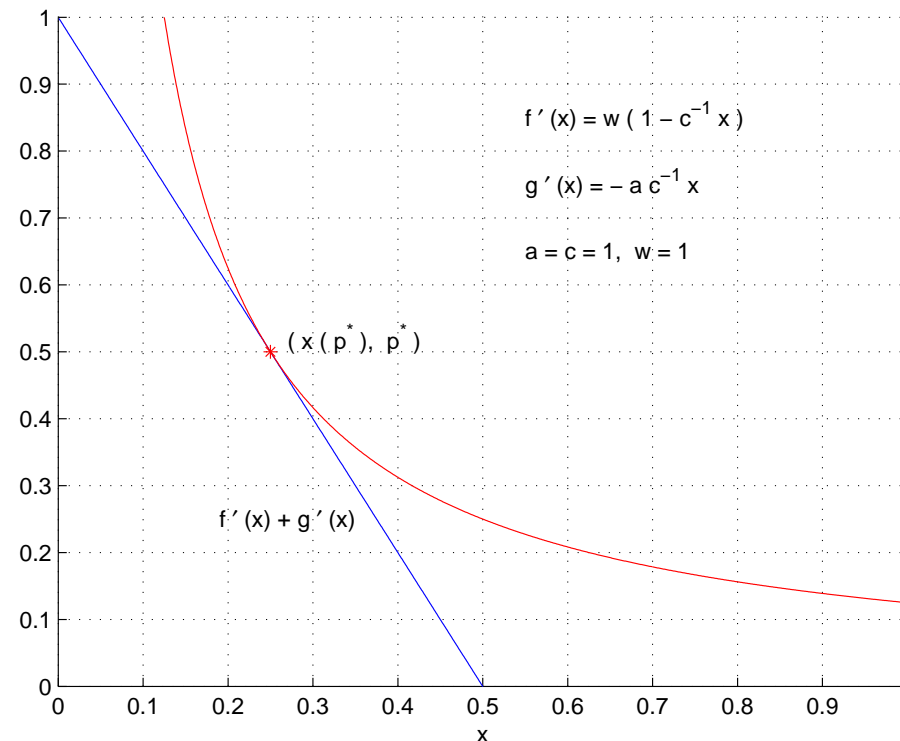


## Single User

- No Classification for Incomplete Information
- Complete Information
- Incomplete Information: Two-point Distribution
- Incomplete Information: Continuous Distribution  
— A Numerical Example
- Comparison of Complete Information and Incomplete Information

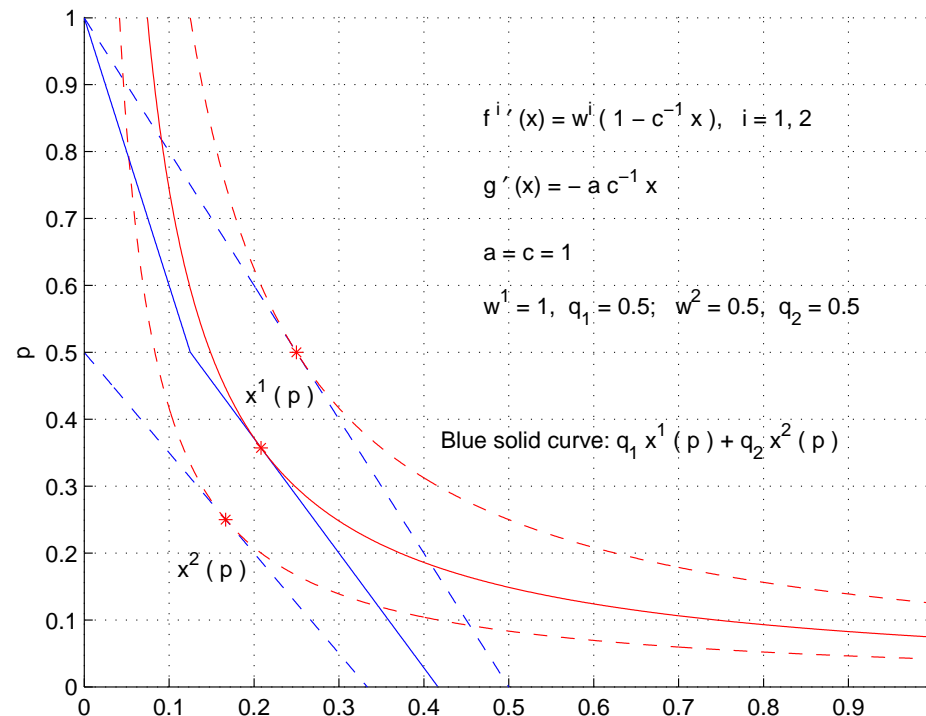
## Complete Information

- User:  $\max_x f(x) + g(x) - px \Rightarrow p = f'(x(p)) + g'(x(p))$
- ISP:  $\max_p px(p)$
- Graphical illustration of the solution



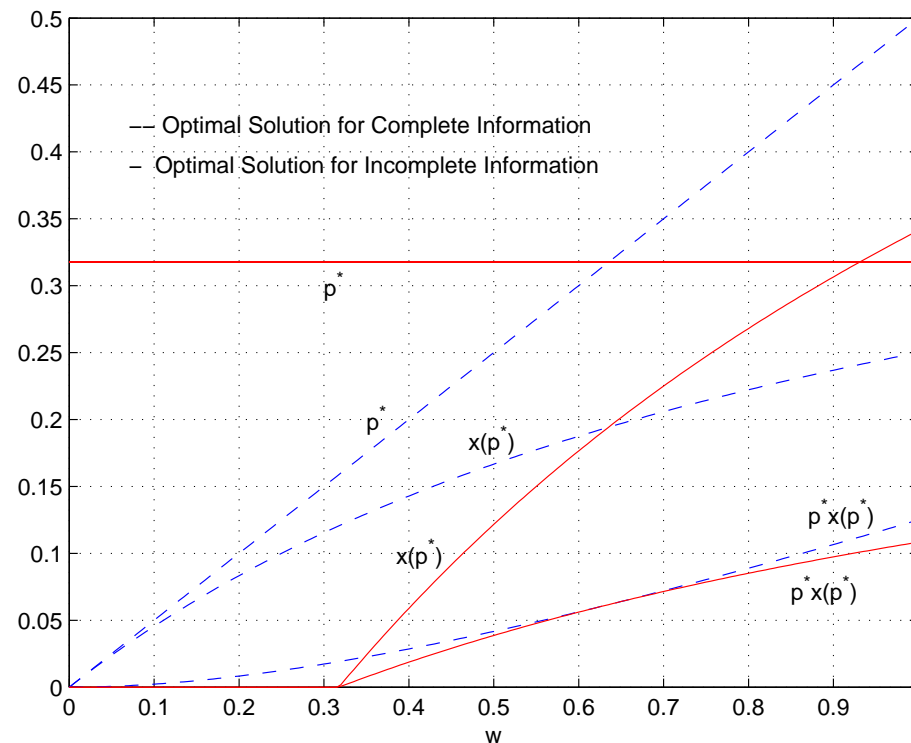
# Incomplete Information — Two-point Distribution

- User:  $x^i(p) = \arg \max_x f^i(x) + g(x) - px$  w.p.  $q_i$ ,  $i = 1, 2$
- ISP:  $p^* = \arg \max_p p[q_1 x^1(p) + q_2 x^2(p)]$
- Complete information:  $p^{i*} = \arg \max_p p x^i(p)$ ,  $i = 1, 2$
- Graphical illustration of the solution and comparison



# Incomplete Information — Uniform Distribution

- $f^{w'}(x) = w(1 - c^{-1}x)$ ,  $g'(x) = -ac^{-1}x$
- $w$  uniformly distributed over  $[0, b]$
- Numerical example:  $a = b = c = 1$
- Profit loss of 12.8% for the ISP due to incomplete information



## Two Users

- Complete Information and Partially Incomplete Information  
— Unique Nash Equilibrium
- Totally Incomplete Information  
— Pure Strategy Bayesian Equilibrium
- Quadratic Utility Functions
  - Complete Information
  - Partially Incomplete Information
  - Totally Incomplete Information
  - Numerical Results
- Comparison of the Three Classes of Games

## Pure Strategy Bayesian Equilibrium

- Special case: independent users  $i = 1, 2$ ,  $f_i(x_i) = f^k(x_i)$  w.p.  $q_k$ ,  $k = 1, 2$
- Pure strategy Bayesian equilibrium:  $x_1^k(\vec{p})$  and  $x_2^k(\vec{p})$ ,  $k = 1, 2$ , solve
 
$$\max_{x_1^k} \{f^k(x_1^k) + q_1 g(x_1^k, x_2^1(\vec{p})) + q_2 g(x_1^k, x_2^2(\vec{p})) - p_1 x_1^k\},$$

$$\max_{x_2^k} \{f^k(x_2^k) + q_1 g(x_1^1(\vec{p}), x_2^k) + q_2 g(x_1^2(\vec{p}), x_2^k) - p_2 x_2^k\}$$
- Sufficient condition for the existence, uniqueness and stability
  - ❶ apply the techniques in Li and Başar (1987)
  - ❷ convergence by the Banach contraction mapping theorem

---

S. Li and T. Başar (1987), “Distributed algorithms for the computation of noncooperative equilibria,” *Automatica*, 23: 523–533.

## Quadratic Utility Functions

- $f_i^{w_i'}(x_i) = w_i(1 - c^{-1}x_i), i = 1, 2$   
 $\nabla g(\vec{x}) = -ac^{-1}(x_1 + x_2)(1, 1)^T$   
independent users:  $i = 1, 2, w_i = w^k$  w.p.  $q_k, k = 1, 2$
- Analytical results for the three types of games
- Numerical results:  
 $a = c = 1, w^1 = 2, w^2 = 1, q_1 = q_2 = 0.5$

# Numerical Results and Comparison of Games

$(w_1, w_2)$	Game	Optimal Price Vector	Optimal Flow Vector	ISP's Profit
(2, 2)	C	(1, 1)	(0.25, 0.25)	0.5
	P	(0.3, 0.3)	(0.4250, 0.4250)	0.2550
	T	(0.7, 0.7)	(0.3647, 0.3647)	0.5106
(2, 1)	C	(1, 0.5)	(0.3, 0.1)	0.35
	P	(0.3, 0.3)	(0.5400, 0.0800)	0.1860
	T	(0.7, 0.7)	(0.3647, 0.0471)	0.2882
(1, 2)	C	(0.5, 1)	(0.1, 0.3)	0.35
	P	(0.3, 0.3)	(0.0800, 0.5400)	0.1860
	T	(0.7, 0.7)	(0.0471, 0.3647)	0.2882
(1, 1)	C	(0.5, 0.5)	(0.1667, 0.1667)	0.1667
	P	(0.3, 0.3)	(0.2333, 0.2333)	0.1400
	T	(0.7, 0.7)	(0.0471, 0.0471)	0.0659

Expected Profit: C — 0.3417   P — 0.1918   T — 0.2882



## Multiple Users

- Quadratic Utility Functions
  - Complete Information
  - Partially Incomplete Information
  - Totally Incomplete Information
  - Numerical Results
- Comparison of the Three Classes of Games

## Quadratic Utility Functions

- $f_i^{w_i'}(x_i) = w_i(1 - x_i)$ ,  $\frac{\partial g(\vec{x})}{\partial x_i} = -\sum_{j=1}^I x_j$   
Independent users:  $1 \leq i \leq I$
- Complete Information: Nash equilibrium — extension of differentiated pricing in Shen and Başar (2004)
- Partially Incomplete Information: Nash equilibrium — extension of uniform pricing in Shen and Başar (2004)
- Totally Incomplete Information: analytical results for the special case with four user types

---

H.-X. Shen and T. Başar (2004), “Differentiated Internet pricing using a hierarchical network game model,” *Proc. IEEE ACC 2004*, pp. 2322–2327.

# Numerical Results and Comparison of Games

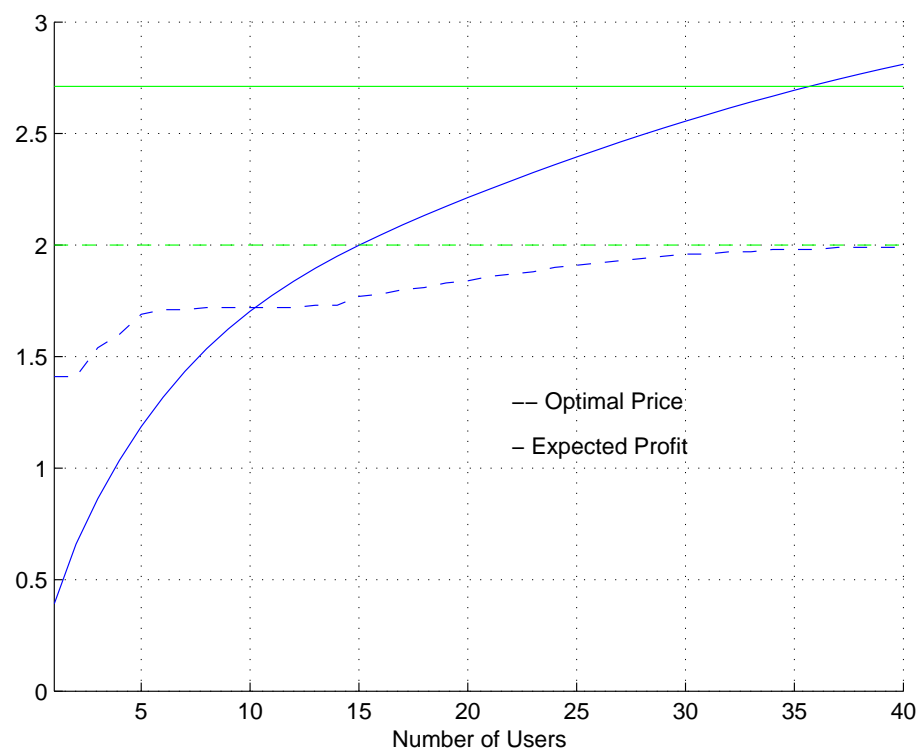
➤ Numerical results:

$w^t = t$  w.p.  $q_t = \frac{1}{4}$ ,  $t \in \{1, 2, 3, 4\}$ ;  $I = 40, 10$  for each type

	C		P		T	
$i, w_i$	$p_i^*$	$x_i(\bar{p}^*)$	$p^*$	$x_i^{\vec{w}}(p^*)$	$p^*$	$x^t(p^*)$
1–10, 4	2	0.1341	1.99	0.1436	2	0.1356
11–20, 3	1.5	0.0122	1.99	0	2	0
21–30, 2	$\geq 0.5366$	0	1.99	0	2	0
31–40, 1	$\geq 0$	0	1.99	0	2	0
$n$	20		10		10	
$\sum x_j$	1.4634		1.4357		1.3559	
$\sum p_j x_j$	2.8659		2.8571		2.7119	

## Partially? Totally?

➤ Partially Incomplete Information vs Totally Incomplete Information



## Conclusions

- Formulation of the three classes of games
- Complete Information vs Incomplete Information:  
Complete Information — ISP and less aggressive users  
Incomplete Information — more aggressive users
- Partially Incomplete Information vs Totally Incomplete Information for the ISP:  
Partially Incomplete Information — with a large number of users  
Totally Incomplete Information — with a small number of users

## Extensions

- General utility functions; Multiple ISPs; . . .

\_\_\_\_\_ End of the Talk \_\_\_\_\_