Hierarchical Network Games with Various Types of Public and Private Information

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Motivation

- Hierarchical network game
 - Stackelberg (leader-follower) game single service provider — leader multiple users — followers
 - Noncooperative game among the users
- A user's type:
 - public information (shared with all the players)
 → Complete information
 - private information (only to that user)
 → Totally incomplete information
 - common information only among the users
 → Partially incomplete information

 Goal: for the three classes of games, study equilibria and compare performances of the players

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Previous Work

- Başar and Srikant INFOCOM'02, JOTA'02 proposed the hierarchical Stackelberg game model studied uniform pricing
- Shen and Başar ACC'04 studied differentiated pricing
- Shen and Başar CDC'04 extended complete information to incomplete information
- This work: extends to general utility functions general distributions for user characteristics





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- Oligie Equilibrium
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- 5 Asymptotic Case



Problem Formulation

- Set of users: $N = \{1, \dots, n\}$; flow of User *i*: x_i , $i \in N$
- Price per unit capacity charged to User i: p_i
- User *i*'s net utility: $f(x_i; w_i) g(x_1, \dots, x_n) p_i x_i$
- Service provider's revenue: $\sum_{j \in N} p_j x_j$
- Users' types (*w_i*'s): independently and identically distributed common distribution publicly known to all the players
- Notations: $\mathbf{w} := \{w_j\}_{j \in N}, \mathbf{x} := \{x_j\}_{j \in N}, \mathbf{p} := \{p_j\}_{j \in N}, \mathbf{w}_{-i} := \{w_j\}_{j \neq i, j \in N}, \mathbf{x}_{-i} := \{x_j\}_{j \neq i, j \in N}.$

Complete Information

Nash equilibrium

Given p:

$$\mathbf{x}_i^c(\mathbf{w};\mathbf{p}) = \arg \max_{\mathbf{x}_i} \{f(\mathbf{x}_i;\mathbf{w}_i) - g(\mathbf{x}_i,\mathbf{x}_{-i}^c(\mathbf{w};\mathbf{p})) - p_i \mathbf{x}_i\}, \ i \in N.$$

Stackelberg game solution

$$\mathbf{p}^{c}(\mathbf{w}) = \arg \max_{\mathbf{p}} \sum_{j \in N} p_{j} x_{j}^{c}(\mathbf{w}; \mathbf{p}), \ \mathbf{x}^{c}(\mathbf{w}; \mathbf{p}^{c}(\mathbf{w})).$$

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Partially Incomplete Information

Nash equilibrium

Given $p_i = p, i \in N$:

$$\mathbf{x}_{i}^{p}(\mathbf{w};\boldsymbol{p}) = \arg \max_{\mathbf{x}_{i}} \{f(\mathbf{x}_{i};\mathbf{w}_{i}) - g(\mathbf{x}_{i},\mathbf{x}_{-i}^{p}(\mathbf{w};\boldsymbol{p})) - p\mathbf{x}_{i}\}.$$

Stackelberg game solution

$$p^{\rho} = \arg \max_{\rho} \rho E_{\mathbf{w}}[\sum_{j \in N} x_j^{\rho}(\mathbf{w}; \rho)], \ \mathbf{x}^{\rho}(\mathbf{w}; \rho^{\rho}).$$

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Totally Incomplete Information

Bayesian equilibrium

Given $p_i = p, i \in N$:

$$\mathbf{x}_i^t(\mathbf{w}_i; \mathbf{p}) = \arg \max_{\mathbf{x}_i} \{ f(\mathbf{x}_i; \mathbf{w}_i) - \mathbf{E}_{\mathbf{w}_{-i}}[g(\mathbf{x}_i, \mathbf{x}_{-i}^t(\mathbf{w}_{-i}; \mathbf{p}))] - \mathbf{p}\mathbf{x}_i \}.$$

Stackelberg game solution

$$\boldsymbol{\rho}^t = \arg\max_{\boldsymbol{\rho}} \boldsymbol{\rho} \boldsymbol{n} \boldsymbol{E}_{\boldsymbol{w}_i}[\boldsymbol{x}_i^t(\boldsymbol{w}_i;\boldsymbol{\rho})], \ \{\boldsymbol{x}_i^t(\boldsymbol{w}_i;\boldsymbol{\rho}^t)\}_{i\in N}.$$

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Existence of a Unique Nash Equilibrium

Theorem

For the complete information or partially incomplete information game, given the prices, there exists a unique Nash equilibrium.

Proof

1. For User $i, i \in N$, the following two are equivalent:

$$\max_{x_i} f(x_i; w_i) - g(x_1, \cdots, x_n) - p_i x_i,$$

$$\max_{x_i}\sum_{j\in N}f(x_j;w_j)-g(x_1,\cdots,x_n)-\sum_{j\in N}p_jx_j.$$

2. The common objective function is strictly concave, and the unique maximizing *n*-tuple is the unique Nash equilibrium.

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Existence of a Unique Bayesian Equilibrium

Theorem

For the totally incomplete information game, given the prices, there exists a unique Bayesian equilibrium.

Proof

Assume $w_i = w^l$ w.p. q_l , $l \in M$. For User *i* with the type w^{l_i} , $i \in N$, $l_i \in M$, the following two are equivalent:

$$\max_{x_{i}^{l_{i}}} f(x_{i}^{l_{i}}; w^{l_{i}}) - \sum_{\{l_{j}\}_{j \neq i}, l_{j} \in M} (\prod_{j \neq i} q_{l_{j}}) g(x_{1}^{l_{1}}, \cdots, x_{n}^{l_{n}}) - p x_{i}^{l_{i}},$$

 $\max_{x_i^{l_j}} \sum_{j \in N} \sum_{l_j \in M} q_{l_j}[f(x_j^{l_j}; w^{l_j}) - px_j^{l_j}] - \sum_{\{l_j\}_{j \in N, l_j \in M}} (\prod_{j \in N} q_{l_j})g(x_1^{l_1}, \cdots, x_n^{l_n}).$

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Two Users with the Same True Types

Three sets of types: \underline{S} , S_0 and \overline{S} (sup $\underline{S} < w_0 \in S_0 < \inf \overline{S}$)

$$egin{aligned} & w \in \overline{S} & \Rightarrow & x_i^{\mathcal{P}}(\mathbf{w}; \mathcal{p}) < x_i^t(w; \mathcal{p}) \ & w \in S_0 & \Rightarrow & x_i^{\mathcal{P}}(\mathbf{w}; \mathcal{p}) = x_i^t(w; \mathcal{p}) \ & w \in \underline{S} & \Rightarrow & x_i^{\mathcal{P}}(\mathbf{w}; \mathcal{p}) > x_i^t(w; \mathcal{p}) \end{aligned}$$

Partially vs Totally incomplete information

 $\begin{array}{l} \text{high types} \Rightarrow \\ \text{more aggressive users under totally incomplete information} \\ \text{low types} \Rightarrow \\ \text{more aggressive users under partially incomplete information} \end{array}$

Two Users with General True Types

$$w_1 \in \overline{S}$$
 $w_1 \leq w_2$ \Rightarrow $x_1^p(\mathbf{w}; p) < x_1^t(w_1; p)$

$$egin{aligned} & w_1 < w_2 & \Rightarrow & x_1^{\mathcal{P}}(\mathbf{w}; \mathcal{p}) \leq x_1^t(w_1; \mathcal{p}) \ w_1 \in S_0 & w_1 = w_2 & \Rightarrow & x_1^{\mathcal{P}}(\mathbf{w}; \mathcal{p}) = x_1^t(w_1; \mathcal{p}) \ w_1 > w_2 & \Rightarrow & x_1^{\mathcal{P}}(\mathbf{w}; \mathcal{p}) \geq x_1^t(w_1; \mathcal{p}) \end{aligned}$$

$$w_1 \in \underline{S}$$
 $w_1 \ge w_2$ \Rightarrow $x_1^{p}(\mathbf{w}; p) > x_1^{t}(w_1; p)$

Partially vs Totally incomplete information

high type User 1, higher type User $2 \Rightarrow$ more aggressive User 1 under totally incomplete information low type User 1, lower type User 2 \Rightarrow more aggressive User 1 under partially incomplete information

Service Provider's Game Preference

Conclusions

- Intuitively, high type users tend to act more aggressively under totally incomplete information, which leads to a higher profit for the service provider.
- On the other hand, when users have comparatively low types, the service provider may prefer the partially incomplete information game for a higher profit.
- For a small number of users, precise comparison results are hard to obtain, even with a fixed pricing policy.

Asymptotic Case

- Communication networks: a large number of users
- Wardrop equilibrium
 - transportation networks neglect an individual vehicle's impact on the total traffic
 - networks with a large number of users
 Nash equilibrium → Wardrop equilibrium (Haurie and Marcotte '85)
 - communication networks routing and flow control (Altman, Başar and Srikant '02, Acemoglu and Ozdaglar '06)

Wardrop Equilibrium

- As $n \to \infty$, for User *i*, regard $g(x_1, \dots, x_n)$ as a constant.
- Given the prices, revise the users' game as follows: complete information: max_{xi} f(x_i; w_i) − p_ix_i, i ∈ N; incomplete information: max_{xi} f(x_i; w_i) − px_i, i ∈ N.
- For a special class of utility functions (as in Başar and Srikant INFOCOM'02):

$$f(x_i; w_i) = w_i \log(1 + x_i), \quad i \in N,$$

$$g(x_1, \cdots, x_n) = \frac{1}{nc - \sum_{j \in N} x_j},$$

the asymptotic behavior of the Nash / Bayesian equilibrium converges to the Wardrop equilibrium.

Service Provider's Game Preference

For the special class of utility functions:

Conclusions

- The service provider makes the highest profit under complete information, since he can charge differentiated prices according to the users' true types.
- The service provider makes a higher expected profit under partially incomplete information; under totally incomplete information, he charges more conservatively (with a higher optimal price) to guarantee that the total capacity is not exceeded.

Extensions

- Asymptotic case: convergence to Wardrop equilibrium for general utility functions
- Nonlinear pricing (with quantity discount)
- Multiple service providers